# BINARY SEQUENCES WITHOUT $0\underbrace{11\dots11}_{k-1}0$ FOR FIXED k

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Abstract. The paper gives a special construction of those words (binary sequences) of length n over alphabet  $\{0,1\}$  in which the subword  $0 \underbrace{11 \dots 11}_{k-1} 0$  is forbidden for some natural number k.

This number of words is counted in two different ways, which gives some new combinatorial identities.

## 1. Definitions and notations

Let  $X = \{0, 1\}$  denote 2-element set of digits. X is called an alphabet. By  $X^n$  we shall denote the set of all strings of length n over alphabet X, i.e.

$$X^{n} = \{ x_{1}x_{2} \dots x_{n} \mid x_{1} \in X \land x_{2} \in X \land \dots \land x_{n} \in X \},\$$

the only element of  $X^0$  is the empty string, i.e. the string of the length 0. The set of all finite strings over alphabet X is

$$X^* = \bigcup_{n \ge 0} X^n.$$

If S is a set, then |S| is the cardinality of S. By  $\lceil x \rceil$  and  $\lfloor x \rfloor$  we denote the smallest integer  $\geq x$  and the greatest integer  $\leq x$ , respectively. By  $\ell_0(p)$  and  $\ell_1(p)$  we denote the number of zeros and ones respectively in the string  $p \in X^*$ .  $N_n = \{1, 2, \ldots, n\}, N_n = \emptyset$  for  $n \leq 0$ ,  $\binom{n}{k} = 0$  iff n < k and [x] is the nearest integer to x.

## 2. Results and discusion

Now we shall construct and enumerate the set of words

$$A_k(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-k}) (x_i x_{i+1} \dots x_{i+k} \neq 0 \underbrace{1 \dots 1}_{k-1} 0) \}$$

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for each natural number k. It is known that

$$a_1(n) = |A_1(n)| = \sum_{i=0}^{\lceil n/2 \rceil} \binom{n-i+1}{i} = \left[\frac{5+3\sqrt{5}}{10} \left(\frac{1+\sqrt{5}}{2}\right)^n\right]$$
(1)

(Fibonacci numbers) where

$$A_1(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-1}) (x_i x_{i+1} \neq 00) \}.$$

In [5] it is shown that the following theorem is valid.

THEOREM 1.

$$a_2(n) = |A_2(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{i-1} \binom{i-1}{j} \binom{n-i-j+1}{j+1} = \left[\frac{2\alpha^2+1}{2\alpha^2-2\alpha+3}\alpha^n\right]$$

where

$$A_2(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-2}) (x_i x_{i+1} x_{i+2} \neq 010) \}$$

and

$$\alpha = \frac{1}{6} \left(4 + \sqrt[3]{100 + 4\sqrt{621}} + \sqrt[3]{100 - 4\sqrt{621}}\right) \approx 1,754877666247.$$

 $A_2(n)$  is the set of all words of length n over alphabet  $\{0,1\}$  with forbidden subword 010.

THEOREM 2.

$$|A_3(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{3} \rfloor} {i-1 \choose j} {i-1-j \choose k} {n-i-j-2k+1 \choose k+1}$$

where

$$A_3(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-3}) (x_i x_{i+1} x_{i+2} x_{i+3} \neq 0110) \}.$$

*Proof.* Now we shall construct this set of words  $A_3(n)$  in some special way, which gives the result for  $|A_3(n)|$ . We make a partition of the set  $A_3(n)$  into subsets  $A_3^i(n)$  which contain exactly *i* zeros.

$$A_3^i(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} \in A_3(n), (\forall s \in N_{n-3}) (x_s x_{s+1} x_{s+2} x_{s+3} \neq 0110), \ell_0(\mathbf{x_n}) = i \}.$$

First we write *i* zeros and then we write one of the letters from the set  $\{p, q, \lambda\}$ on the i - 1  $(1 \le i \le n)$  places between *i* zeros where p = 1, q = 111 and  $\lambda$  is the empty letter. Letter  $\lambda$  is the letter with property that if  $\lambda$  is writen between two zeros then actually nothing is written. Let *j* be the number of appearances of the letter *p* and *k* is the number of appearances of the letter *q*. We choose *j* places from

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i-1 places for letters p and after that we choose k places from i-1-j places for letters q. This we can do in

$$\binom{i-1}{j}\binom{i-1-j}{k} \tag{2}$$

different ways. Now we have only n-i-j-3k ones, which must be put on k places where we have subwords 111 as well as into the regions in front of and behind the word, that is into k + 2 regions in all. It can be done in

$$\binom{n-i-j-2k+1}{k+1} \quad \text{ways} . \tag{3}$$

Thus from (2), (3) and

$$|A_3(n)| = \sum_{i=0}^n |A_3^i(n)| = 1 + \sum_{i=1}^n |A_3^i(n)|$$

Theorem 2 follows.  $\blacksquare$ 

THEOREM 3.

$$a_{3}(n) = |A_{3}(n)| = \left[\frac{2\alpha^{3} + 1}{2\alpha^{3} - 3\alpha + 4}\alpha^{n}\right]$$
  
$$\alpha = \frac{1}{2\alpha^{3} - 3\alpha + 4}\alpha^{n} = \frac{1}{2\alpha^{3} - 3\alpha + 4}\alpha^{n}$$

where

$$\alpha = \frac{1}{2} \left( 1 + \sqrt{3 + 2\sqrt{5}} \right) \approx 1,866760399.$$

*Proof.* Words  $\mathbf{x_n} \in A_3(n)$  are obtained from other words  $\mathbf{x_{n-1}} \in A_3(n-1)$  by appending 0 or 1 in front of them. Let  $\mathbf{x_{n-1}} \in A_3(n-1)$ ,  $\mathbf{x_{n-3}} \in A_3(n-3)$  and  $\mathbf{x_{n-4}} \in A_3(n-4)$ . Then  $1\mathbf{x_{n-1}} \in A_3(n)$ ,  $0110\mathbf{x_{n-4}} \notin A_3(n)$  and  $0111\mathbf{x_{n-4}} \in A_3(n)$ , which means that  $011\mathbf{x_{n-3}} \in A_3(n)$  if and only if  $\mathbf{x_{n-3}}$  begins with letter 1. This implies the recourse relation

$$a_3(n) = 2a_3(n-1) - a_3(n-3) + a_3(n-4)$$

whose characteristic equation is  $x^4 - 2x^3 + x - 1 = 0$  and whose roots are

$$\alpha = \frac{1}{2} \left( 1 + \sqrt{3 + 2\sqrt{5}} \right), \quad \beta = \frac{1}{2} \left( 1 - \sqrt{3 + 2\sqrt{5}} \right)$$
$$\gamma = \frac{1}{2} \left( 1 + i\sqrt{2\sqrt{5} - 3} \right) \quad \text{and} \quad \delta = \frac{1}{2} \left( 1 - i\sqrt{2\sqrt{5} - 3} \right).$$

The explicit formula for  $a_3(n)$  is

$$a_3(n) = C_1 \alpha^n + C_2 \beta^n + C_3 \gamma^n + C_4 \delta^n$$

where

$$C_1 = \frac{2\alpha^3 + 1}{2\alpha^3 - 3\alpha + 4}, \quad C_2 = \frac{2\beta^3 + 1}{2\beta^3 - 3\beta + 4}$$
$$C_3 = \frac{2\gamma^3 + 1}{2\gamma^3 - 3\gamma + 4}, \quad \text{and} \quad C_4 = \frac{2\delta^3 + 1}{2\delta^3 - 3\delta + 4}.$$

Since  $|\beta| < 1$ ,  $|\gamma| < 1$  and  $|\delta| < 1$  we obtain Theorem 3.

Thus from Theorem 2 and Theorem 3 follows

Corollary 1.

$$|A_{3}(n)| = 1 + \sum_{i=1}^{n} \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{2} \rfloor} {i-1 \choose j} {i-1-j \choose k} {n-i-j-2k+1 \choose k+1} = \left[ \frac{2\alpha^{2}+1}{2\alpha^{3}-3\alpha+4} \alpha^{n} \right], \quad where \quad \alpha = \frac{1}{2} \left( 1 + \sqrt{3+2\sqrt{5}} \right).$$

THEOREM 4.

$$|A_k(n)| = 1 + \sum_{i=1}^n \sum_{j_1=0}^{n-i} \sum_{j_2=0}^{\lfloor \frac{n-i-S_1}{2} \rfloor} \sum_{j_3=0}^{\lfloor \frac{n-i-S_2}{3} \rfloor} \cdots \sum_{j_{k-2}=0}^{\lfloor \frac{n-i-S_{k-2}}{k-2} \rfloor} \sum_{\ell=0}^{\lfloor \frac{n-i-S_{k-2}}{k} \rfloor} \prod_{m=0}^{l-i-S_{k-2}} \left( \frac{i-1-s_m}{j_{m+1}} \right) \binom{i-1-s_{k-2}}{\ell} \binom{n-i-S_{k-2}-(k-1)\ell+1}{\ell+1}$$

where  $s_k = j_1 + j_2 + \dots + j_k$ ,  $s_0 = 0$ ,  $S_k = j_1 + 2j_2 + \dots + kj_k$  and  $A_k(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall s \in N_{n-k}) (x_s x_{s+1} \dots x_{s+k} \neq 0 \underbrace{1 \dots 1}_{k-1} 0) \}.$ 

*Proof.* We partition the set  $A_k(n)$  into subsets  $A_k^i(n)$  which contain exactly *i* zeros i.e.

$$A_{k}^{i}(n) = \{ \mathbf{x_{n}} \mid \mathbf{x_{n}} = x_{1}x_{2}\dots x_{n} \in X^{n}, \\ (\forall s \in N_{n-k})(x_{s}x_{s+1}\dots x_{s+k} \neq 0 \underbrace{1\dots 1}_{k-1} 0), \ell_{0}(\mathbf{x_{n}}) = i \}.$$

Now we shall construct words from  $A_k^i(n)$  in the following way. First we write i zeros and then we write one of the letters from the alphabet  $\{q_1, q_2, \ldots, q_{k-2}, r, \lambda\}$  on i-1 places between i zeros where  $q_m = \underbrace{11 \ldots 1}_m$ , for  $m \in \{1, 2, \ldots, k-2\}$ ,  $r = \underbrace{11 \ldots 1}_k$  and  $\lambda$  is the empty letter. Let  $j_m$  be the number of letters  $q_m$ , and  $\ell$  the number of letters r. We choose  $j_1$  places from i-1 places for letters  $q_1, j_2$  places from  $i-1-j_1$  places for letters  $q_2, \ldots, j_{k-2}$  places from  $i-1-s_{k-3}$  places for letters r. It can be done in

$$\prod_{m=0}^{m=k-3} \binom{i-1-s_m}{j_{m+1}} \binom{i-1-s_{k-2}}{\ell}$$
(4)

different ways, where  $s_k = j_1 + j_2 + \cdots + j_k$  and  $s_0 = 0$ . There remains to write  $n - i - S_{k-2} - k\ell$  letters 1 on  $\ell$  regions which already contain r, as well as into the

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regions in front of and behind the word, that is into  $\ell+2$  regions in all. It can be done in

$$\binom{n-i-S_{k-2}-(k-1)\ell+1}{\ell+1}$$
(5)

different ways, where  $S_k = j_1 + 2j_2 + \cdots + kj_k$ . Thus from (4), (5) and  $|A_k(n)| = \sum_{i=0}^n |A_k^i(n)|$  Theorem 4 follows.

Theorem 5.

$$|A_k(n)| = [C(k,\alpha)\alpha^n]$$

for large enough values of n, where  $\alpha$  is the unique real root of equation

$$x^{k+1} - 2x^k + x - 1 = 0$$

which lies between 1 and 2 and  $C(k, \alpha)$  is the rational function of  $\alpha$  and k.

*Proof.* Words  $\mathbf{x_n} \in A_k(n)$  are obtained from other words  $\mathbf{x_{n-1}} \in A_k(n-1)$  by appending 0 or 1 in front of them. Let

$$\mathbf{x_{n-1}} \in A_k(n-1), \ \mathbf{x_{n-k}} \in A_k(n-k) \text{ and } \mathbf{x_{n-k-1}} \in A_k(n-k-1).$$

Then

$$1\mathbf{x_{n-1}} \in A_k(n), \ 0\underbrace{11\dots 1}_k \mathbf{x_{n-k-1}} \in A_k(n), \ 0\underbrace{11\dots 1}_{k-1} 0\mathbf{x_{n-k-1}} \notin A_k(n)$$

which means that  $0\underbrace{11\ldots 1}_{k-1}\mathbf{x_{n-k}} \in A_k(n)$  if and only if  $\mathbf{x_{n-k}}$  begins with the latter 1. This implies the recommence relation

letter 1. This implies the reccurence relation

$$a_k(n) = 2a_k(n-1) - a_k(n-k) + a_k(n-k-1)$$

whose characteristic equation is  $x^{k+1} - 2x^k + x - 1 = 0$  which has only one real root  $\alpha$  for k = 2m,  $m \in N$ . This real root lies between 1 and 2. If k = 2m + 1,  $m \in N \cup \{0\}$ , then the characteristic equation has only two real roots  $\alpha$  and  $\beta$ where  $\alpha \in (1,2)$  and  $\beta \in (-1,0)$ . The complex roots have modules less than 1. Because of that it follows that  $a_k(n) = [C(k, \alpha)\alpha^n]$  for large enough values of n, where  $C(k, \alpha)$  is rational function of  $\alpha$  and k.

Corollary 2.

$$|A_4(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{4} \rfloor} \sum_{\ell=0}^{\lfloor \frac{n-i-j-2k}{4} \rfloor} \left( \binom{i-1}{j} \binom{i-1-j}{k} \binom{i-1-j-2k}{\ell} \binom{n-i-j-2k-3\ell+1}{\ell+1} \right) \\ = \left[ \frac{4\alpha^4 - \alpha + 2}{4\alpha^4 - 4\alpha^2 + 3\alpha + 2} \alpha^n \right],$$

i.e.  $C(4, \alpha) = \frac{4\alpha^4 - \alpha + 2}{4\alpha^4 - 4\alpha^2 + 3\alpha + 2}$  and  $\alpha$  is unique real root of equation  $x^5 - 2x^4 + x - 1 = 0$  whose complex roots are with modules less than 1.

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