BINARY SEQUENCES WITHOUT 0 11 ... 11 [|] {z } k1 ⁰ FOR FIXED k

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Abstract. The paper gives a special construction of those words (binary sequences) of length *n* over alphabet $\{0,1\}$ in which the subword $0\underbrace{11\ldots11}_{k-1}0$ is forbidden for some natural number *k*. [|] {z }

This number of words is counted in two different ways, which gives some new combinatorial identities.

1. Definitions and notations

Let $X = \{0, 1\}$ denote 2-element set of digits. X is called an alphabet. By X^n we shall denote the set of all strings of length ν over alphabet Λ , i.e. 1. **Definitions and notations**
{0,1} denote 2-element set of digits. *X* is called an alphabet. By *X* is the set of all strings of length *n* over alphabet *X*, i.e.
 $X^n = \{x_1x_2...x_n | x_1 \in X \land x_2 \in X \land \cdots \land x_n \in X\}$,

$$
X^n = \{ x_1 x_2 \dots x_n \mid x_1 \in X \land x_2 \in X \land \dots \land x_n \in X \},
$$

the only element of Λ is the empty string, i.e. the string of the length 0. The set of all nite strings over alphabet X is

$$
X^* = \bigcup_{n \ge 0} X^n.
$$

If S is a set, then $|S|$ is the cardinality of S. By [x] and $|x|$ we denote the smallest integer $\geq x$ and the greatest integer $\leq x$, respectively. By $\ell_0(p)$ and $\ell_1(p)$ we denote the number of zeros and ones respectively in the string $p \in X^*$. $N_n = \{1, 2, \ldots, n\},\ N_n = \emptyset \text{ for } n \leq 0,\ \binom{n}{k} = 0 \text{ iff } n < k \text{ and } [x] \text{ is the nearest }$ integer to x .

2. Results and discusion

Now we shall construct and enumerate the set of words

$$
A_k(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, \ (\forall i \in N_{n-k})(x_i x_{i+1} \dots x_{i+k} \neq 0 \underbrace{1 \dots 1}_{k-1} 0) \}
$$

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for each natural number k. It is known that

$$
a_1(n) = |A_1(n)| = \sum_{i=0}^{\lceil n/2 \rceil} {n-i+1 \choose i} = \left[\frac{5+3\sqrt{5}}{10} \left(\frac{1+\sqrt{5}}{2} \right)^n \right] \tag{1}
$$

(Fibonacci numbers) where

$$
A_1(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-1}) (x_i x_{i+1} \neq 00) \}.
$$

In [5] it is shown that the following theorem is valid.

THEOREM 1.

$$
a_2(n) = |A_2(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{i-1} {i-1 \choose j} {n-i-j+1 \choose j+1} = \left[\frac{2\alpha^2+1}{2\alpha^2-2\alpha+3}\alpha^n\right]
$$

where

$$
A_2(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-2}) (x_i x_{i+1} x_{i+2} \neq 010) \}
$$

and

$$
\alpha = \frac{1}{6}(4 + \sqrt[3]{100 + 4\sqrt{621}} + \sqrt[3]{100 - 4\sqrt{621}}) \approx 1,754877666247.
$$

 $A_2(n)$ is the set of all words of length n over alphabet $\{0,1\}$ with forbidden subword 010.

THEOREM 2.

$$
|A_3(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{3} \rfloor} {i-1 \choose j} {i-1-j \choose k} {n-i-j-2k+1 \choose k+1}
$$

where

$$
A_3(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall i \in N_{n-3})(x_i x_{i+1} x_{i+2} x_{i+3} \neq 0110) \}.
$$

Proof. Now we shall construct this set of words A3(n) in some special way, which gives the result for $|A_3(n)|$. We make a partition of the set $A_3(n)$ into subsets $A_3(n)$ which contain exatly i zeros.

$$
A_3^{\iota}(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} \in A_3(n), (\forall s \in N_{n-3}) (x_s x_{s+1} x_{s+2} x_{s+3} \neq 0110), \ell_0(\mathbf{x_n}) = i \}.
$$

First we write i zeros and then we write one of the letters from the set $\{p, q, \lambda\}$ on the $i-1$ $(1 \le i \le n)$ places between i zeros where $p = 1$, $q = 111$ and λ is the empty letter. Letter is the letter with property that if is writen between two zeros then actually nothing is written. Let j be the number of appearances of the letter p and k is the number of appearances of the letter q. We choose j places from

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$$
0 \overbrace{11...11}^{k-1} 0
$$
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 $i-1$ places for letters p and after that we choose k places from $i-1-j$ places for leters q_i - This we can do in

$$
\binom{i-1}{j}\binom{i-1-j}{k}
$$
 (2)

different ways. Now we have only $n-i-j-3k$ ones, which must be put on k places where we have subwords 111 as well as into the regions in front of and behind theword, that is into $\kappa \pm 2$ regions in all. It can be done in

$$
\binom{n-i-j-2k+1}{k+1} \quad \text{ways} \tag{3}
$$

Thus from (2), (3) and

$$
|A_3(n)| = \sum_{i=0}^{n} |A_3^i(n)| = 1 + \sum_{i=1}^{n} |A_3^i(n)|
$$

Theorem 2 follows. ■

THEOREM 3.

$$
a_3(n) = |A_3(n)| = \left[\frac{2\alpha^3 + 1}{2\alpha^3 - 3\alpha + 4}\alpha^n\right]
$$

$$
a_3(n) = \frac{1}{2}\left(1 + \sqrt{3 + 2\sqrt{5}}\right) \approx 1.866760399
$$

where

$$
\alpha = \frac{1}{2} \left(1 + \sqrt{3 + 2\sqrt{5}} \right) \approx 1,866760399.
$$

Proof. Words $x_n \in A_3(n)$ are obtained from other words $x_{n-1} \in A_3(n-1)$ by appending 0 or 1 in front of them. Let $\mathbf{x}_{n-1} \in A_3(n-1)$, $\mathbf{x}_{n-3} \in A_3(n-3)$ and $x_{n-4} \in A_3(n-4)$. Then $1x_{n-1} \in A_3(n)$, $0110x_{n-4} \notin A_3(n)$ and $0111x_{n-4} \in$ $A_3(n)$, which means that $011\mathbf{x}_{n-3} \in A_3(n)$ if and only if \mathbf{x}_{n-3} begins with letter 1. This implies the reccurence relation

$$
a_3(n) = 2a_3(n-1) - a_3(n-3) + a_3(n-4)
$$

whose characteristic equation is $x^4 - 2x^3 + x - 1 = 0$ and whose roots are

$$
\alpha = \frac{1}{2} \left(1 + \sqrt{3 + 2\sqrt{5}} \right), \quad \beta = \frac{1}{2} \left(1 - \sqrt{3 + 2\sqrt{5}} \right)
$$

$$
\gamma = \frac{1}{2} \left(1 + i\sqrt{2\sqrt{5} - 3} \right) \quad \text{and} \quad \delta = \frac{1}{2} \left(1 - i\sqrt{2\sqrt{5} - 3} \right).
$$

The explicit formula for $u_3(n)$ is

$$
a_3(n) = C_1 \alpha^n + C_2 \beta^n + C_3 \gamma^n + C_4 \delta^n
$$

where

$$
C_1 = \frac{2\alpha^3 + 1}{2\alpha^3 - 3\alpha + 4}, \quad C_2 = \frac{2\beta^3 + 1}{2\beta^3 - 3\beta + 4}
$$

$$
C_3 = \frac{2\gamma^3 + 1}{2\gamma^3 - 3\gamma + 4}, \quad \text{and} \quad C_4 = \frac{2\delta^3 + 1}{2\delta^3 - 3\delta + 4}.
$$

Since $|\beta|$ < 1, $|\gamma|$ < 1 and $|\delta|$ < 1 we obtain Theorem 3.

Thus from Theorem 2 and Theorem 3 follows

Corollary 1.

$$
|A_3(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{3} \rfloor} {i-1 \choose j} {i-1-j \choose k} {n-i-j-2k+1 \choose k+1}
$$

= $\left[\frac{2\alpha^2 + 1}{2\alpha^3 - 3\alpha + 4} \alpha^n \right]$, where $\alpha = \frac{1}{2} \left(1 + \sqrt{3 + 2\sqrt{5}} \right)$.

THEOREM 4.

$$
|A_{k}(n)| = 1 + \sum_{i=1}^{n} \sum_{j_{1}=0}^{n-i} \sum_{j_{2}=0}^{\lfloor \frac{n-i-5_{1}}{2} \rfloor} \sum_{j_{3}=0}^{\lfloor \frac{n-i-5_{2}}{3} \rfloor} \cdots \sum_{j_{k-2}=0}^{\lfloor \frac{n-i-5_{k-3}}{k-2} \rfloor} \sum_{\ell=0}^{\lfloor \frac{n-i-5_{k-2}}{k} \rfloor} \sum_{\ell=0}^{\lfloor \frac{n-i-5_{k-2}}{k} \rfloor}
$$

$$
\prod_{m=0}^{m=k-3} {i-1-s_{m} \choose j_{m+1}} {i-1-s_{k-2} \choose \ell} {n-i-S_{k-2}-(k-1)\ell+1 \choose \ell+1}
$$

where $s_k = j_1 + j_2 + \cdots + j_k$, $s_0 = 0$, $S_k = j_1 + 2j_2 + \cdots + k j_k$ and $A_k(n) = {\mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, (\forall s \in N_{n-k})(x_s x_{s+1} \dots x_{s+k} \neq 0 \underbrace{1 \dots 1}_{s+k} 0).}$

Proof. We partition the set $A_k(n)$ into subsets $A_k(n)$ which contain exactly i zeros i.e.

$$
A_k^i(n) = \{ \mathbf{x_n} \mid \mathbf{x_n} = x_1 x_2 \dots x_n \in X^n, \quad (\forall s \in N_{n-k})(x_s x_{s+1} \dots x_{s+k} \neq 0 \underbrace{1 \dots 1}_{k-1} 0), \ell_0(\mathbf{x_n}) = i \}.
$$

Now we shall construct words from $A_k^+(n)$ in the following way. First we write i zeros and then we write one of the letters from the alphabet $\{q_1, q_2, \ldots, q_{k-2}, r, \lambda\}$ on $i-1$ places between i zeros where $q_m = \underbrace{11 \ldots 1}_{m}$, for $m \in \{1, 2, \ldots, k-2\}$, $r = \underbrace{11 \ldots 1}$ and λ is the empty letter. Let j_m be the number of letters q_m , and ℓ the number of letters r. We choose j_1 places from $i-1$ places for letters q_1, j_2 places from $i - 1 - j_1$ places for letters q_2, \ldots, j_{k-2} places from $i - 1 - s_{k-3}$ places for leters q_{k-2} and ℓ places from $i - 1 - s_{k-2}$ places for letters r. It can be done in

$$
\prod_{m=0}^{m=k-3} {i-1-s_m \choose j_{m+1}} {i-1-s_{k-2} \choose \ell} \tag{4}
$$

different ways, where $s_k = j_1 + j_2 + \cdots + j_k$ and $s_0 = 0$. There remains to write $n - i - S_{k-2} - k\ell$ letters 1 on ℓ regions which already contain r, as well as into the

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regions in front of and behind the word, that is into $\epsilon \pm$ 2 regions in all. It can be done in

$$
\binom{n-i-S_{k-2}-(k-1)\ell+1}{\ell+1} \tag{5}
$$

different ways, where $S_k = j_1 + 2j_2 + \cdots + kj_k$. Thus from (4), (5) and $|A_k(n)| =$ $\sum_{i=0}^{n} |A_k^i(n)|$ Theorem 4 follows.

THEOREM 5.

$$
|A_k(n)| = [C(k, \alpha)\alpha^n]
$$

for large enough values of n, where is the unique real root of equation

$$
x^{k+1} - 2x^k + x - 1 = 0
$$

which lies betwen 1 and 2 and \cup (k; α) is the rational function of α and k.

Proof. Words $\mathbf{x_n} \in A_k(n)$ are obtained from other words $\mathbf{x_{n-1}} \in A_k(n-1)$ by appending 0 or 1 in front of them. Let

$$
x_{n-1} \in A_k(n-1), x_{n-k} \in A_k(n-k)
$$
 and $x_{n-k-1} \in A_k(n-k-1)$.

Then

$$
1\mathbf{x}_{n-1}\in A_k(n),\ 0\underbrace{11\ldots1}_{k}\mathbf{x}_{n-k-1}\in A_k(n),\ 0\underbrace{11\ldots1}_{k-1}0\mathbf{x}_{n-k-1}\notin A_k(n)
$$

which incans that $\cup_{i=1}^{n}$, λ_{n-k} \ldots $\mathbf{1}_{\mathbf{x}_{n-k}} \in A_k(n)$ if and only if \mathbf{x}_{n-k} begins with the $k-1$

letter 1. This implies the reccurence relation

$$
a_k(n) = 2a_k(n-1) - a_k(n-k) + a_k(n-k-1)
$$

whose characteristic equation is $x^{k+1} - 2x^k + x - 1 = 0$ which has only one real root α for $k = 2m$, $m \in N$. This real root lies between 1 and 2. If $k = 2m + 1$, $a_k(n) = 2a_k(n-1) - a_k(n-k) + a_k(n-k-1)$
whose characteristic equation is $x^{k+1} - 2x^k + x - 1 = 0$ which has only one real
root α for $k = 2m$, $m \in N$. This real root lies between 1 and 2. If $k = 2m + 1$,
 $m \in N \cup \{0\}$, then the char where $\alpha \in (1,2)$ and $\beta \in (-1,0)$. The complex roots have modules less than 1. Because of that it follows that $a_k(n) = |C(\kappa, \alpha)\alpha^{\gamma}|$ for large enough values of n , where $C(\kappa, \alpha)$ is rational function of α and κ .

Corollary 2.

$$
|A_4(n)| = 1 + \sum_{i=1}^n \sum_{j=0}^{n-i} \sum_{k=0}^{\lfloor \frac{n-i-j}{2} \rfloor} \sum_{\ell=0}^{\lfloor \frac{n-i-j}{4} \rfloor}
$$

$$
{i-1 \choose j} {i-1-j \choose k} {i-1-j-2k \choose \ell} {n-i-j-2k-3\ell+1 \choose \ell+1}
$$

$$
= \left[\frac{4\alpha^4 - \alpha + 2}{4\alpha^4 - 4\alpha^2 + 3\alpha + 2} \alpha^n \right],
$$

i.e. $C(4,\alpha) = \frac{4\alpha^2-4\alpha^2+3\alpha+2}{4\alpha^4-4\alpha^2+3\alpha+2}$ and α is unique real root of equation $x^5-2x^4+x-1=0$ whose complex roots are with modules less than 1.

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